

# Basic (mathematical) knowledge on the structure of voting systems with several levels of approval

Josep Freixas<sup>a</sup>

<sup>a</sup>Department of Mathematics  
Technical University of Catalonia

Summer School, Campione d'Italia  
Game Theory and Voting Systems



# Outline

- 1 Voting rules with abstention or  $(3, 2)$  simple games
  - The model
- 2 Weighted  $(3, 2)$  games
  - Examples
  - Subclasses of weighted VGAs
- 3 Complete  $(3, 2)$  games
- 4 Linking weightedness and completeness
  - Link between weightedness and completeness
  - Games with few voters
- 5 A little bit on (Banzhaf) power: two versions
  - Computation of power indices



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## Tripartitions:

$$(S_1, S_2, S_3) \in 3^N$$

- $S_1$  set of “yes” voters
- $S_2$  set of “abstainers”
- $S_3$  set of “no” voters



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## Monotonicity:

$$(S_1, S_2, S_3) \sqsubseteq (T_1, T_2, T_3)$$

if

$$S_1 \subseteq T_1 \quad \text{and} \quad S_2 \subseteq T_1 \cup T_2$$



A (3, 2) simple game  $(N, V)$  consists of a finite set  $N$  of voters together with a value function  $V : 3^N \rightarrow \{\text{win}, \text{lose}\}$  such that:

- $(N, \emptyset, \emptyset)$  wins,
- $(\emptyset, \emptyset, N)$  loses,
- $S \sqsubseteq T$  and  $V(S) = \text{wins}$  implies  $V(T) = \text{wins}$ .



## Example

Consider the (3, 2) game defined on  $N = \{1, 2, 3\}$  by its set of minimal winning tripartitions

$W^m = \{(12, \emptyset, 3), (1, 23, \emptyset), (23, 1, \emptyset)\}$  (keys and commas omitted for tripartitions).

(123, $\emptyset$ , $\emptyset$ )	winning but non-minimal
(12, 3, $\emptyset$ )	winning but non-minimal
(13, 2, $\emptyset$ )	winning but non-minimal
(1, 23, $\emptyset$ )	minimal winning
(12, $\emptyset$ , 3)	minimal winning
(23, 1, $\emptyset$ )	minimal winning
others	losing

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A (3, 2) simple game  $(N, W)$  is a **weighted (3, 2)** game if there exists a weight function  $w : N \rightarrow \mathbb{R}^3$  and a quota  $q$  such that

$$S \in W \iff w(S) \geq q$$

where  $w(S) = \sum_{p \in S_1} w^+(p) + \sum_{p \in S_3} w^-(p)$  for all

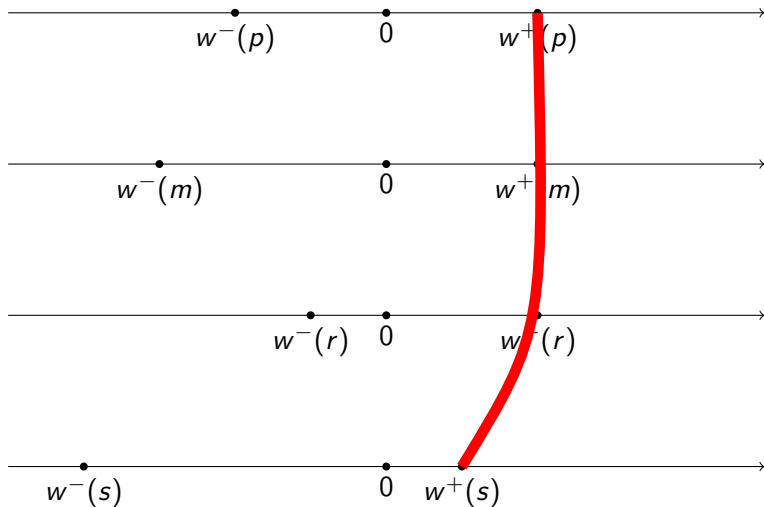
$$S = (S_1, S_2, S_3) \in 3^N.$$

The only requirement for weights is:

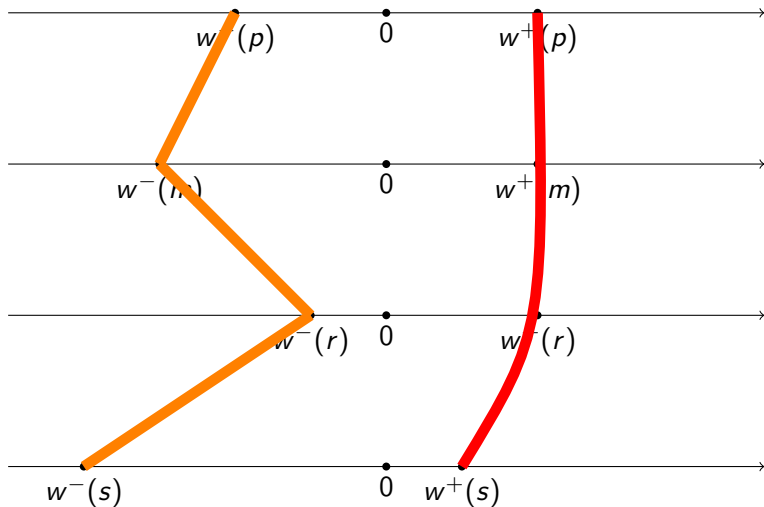
$$w^+(p) \geq 0 \geq w^-(p)$$



## Some weights in a weighted game



## Some weights in a weighted game



## Example

A resolution is carried in the Security Council if at least nine members support it and no permanent member is explicitly opposed. Let  $P = \{1, 2, 3, 4, 5\}$  and  $R = \{6, 7, \dots, 15\}$  be resp. the set of permanent and nonpermanent members, and

$$V(S) = V(S_1, S_2, S_3) = \begin{cases} \text{win} & \text{if } |S_1| \geq 9 \text{ and } S_3 \cap P = \emptyset \\ \text{lose} & \text{otherwise} \end{cases}$$

This voting system with abstention can be represented as

$$[ 9; \underbrace{(1, 0, -6), \dots, (1, 0, -6)}_5, \underbrace{(1, 0, 0), \dots, (1, 0, 0)}_{10} ]$$



## Example

A juror has **three** choices—convict, acquit or be neutral—but the outcome of the jury as a whole has only **two** choices—convict or acquit. Assume there are 4 jurors, A, B, C, and D. The collective decision is “convict” *if and only if*:



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**Version 2** at least one juror votes for “convicting” and at most one juror votes for “acquittal”.

**Version 3** more jurors vote for “convicting” than for “acquittal” and juror A votes for “convicting”.





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Voting rule in **version 1** is weighted, with  $w(p) = (1, 0, -1)$  for each  $p \in N$  and  $q = 1$ .

$$\left[ 1; \underbrace{(1, 0, -1), \dots, (1, 0, -1)}_4 \right]$$

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Voting rule in **version 2** is **not weighted!** (Check it!)

Voting rule in **version 3** is weighted. Two representations:

$$\left[ 4; (4, 0, 0), (1, 0, -1), (1, 0, -1), (1, 0, -1) \right]$$

$$\left[ 4; (4, 0, -1), (1, 0, -1), (1, 0, -1), (1, 0, -1) \right]$$

Two consecutive stronger conditions of a weighted (3, 2):

A **strongly weighted (3, 2) game** is a weighted (3, 2) game that admits a representation such that for every pair of voters  $p$  and  $r$ , either

$$w^+(p) \geq w^+(r) \quad \text{and} \quad -w^-(p) \geq -w^-(r)$$

or

$$w^+(p) \leq w^+(r) \quad \text{and} \quad -w^-(p) \leq -w^-(r).$$



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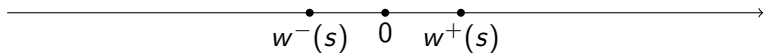
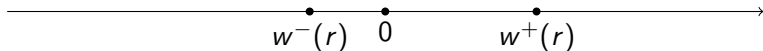
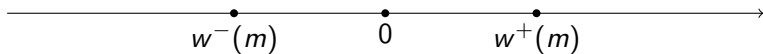
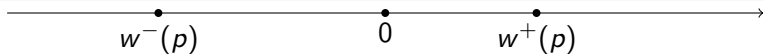
or

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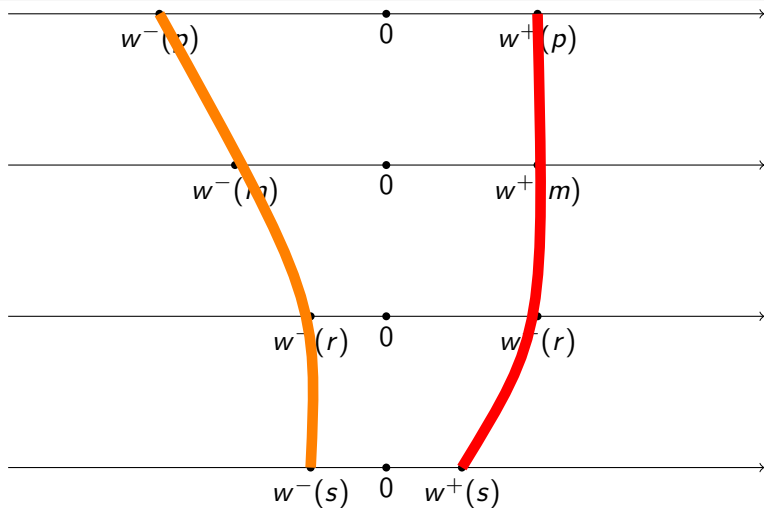
A **zero-centered strongly weighted (3, 2) game** is a strongly weighted (3, 2) game that admits a representation with weights  $w^+(p) = -w^-(p)$  for each voter  $p \in N$ .



## Weights in a strongly weighted game

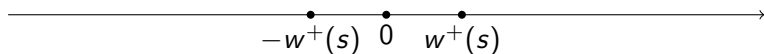
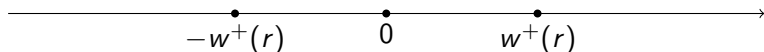
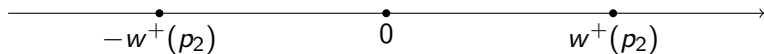
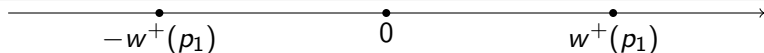


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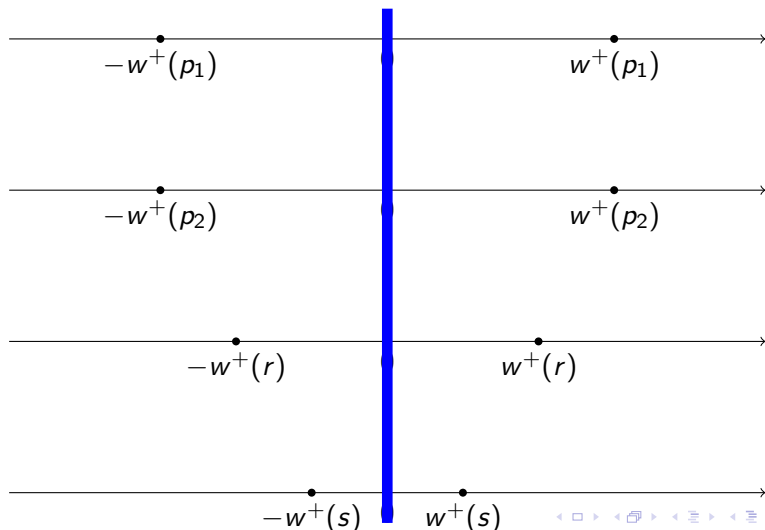




## Weights in a zero-centered strongly weighted game



## Weights in a zero-centered strongly weighted game



## Inclusion of weighted VRA

weighted  
games

strongly  
weighted games

zero-centered  
weighted games

## Examples revisited

- The 4 jurors example in version 1 is a **zero-centered** strongly weighted (3,2) game.  $[ 1; \underbrace{(1, 0, -1), \dots, (1, 0, -1)}_4 ]$



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- The UNSC voting system with abstention is a **strongly** weighted (3,2) game.  
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Tchantcho et. al. (2008). The *I*-influence relation.

Let  $(N, V)$  be a (3, 2) game,  $p, r \in N$ . Voter  $p$  is said to be at least as *influential* as  $r$ ,  $p \succsim_I r$ , if for all  $(S_1, S_2, S_3) \in 3^N$  it yields:

- $V(S_1 \cup p, S_2 \setminus p, S_3) \geq V(S_1 \cup r, S_2 \setminus r, S_3)$  if  $p, r \in S_2$ ,
- $V(S_1, S_2 \cup p, S_3 \setminus p) \geq V(S_1, S_2 \cup r, S_3 \setminus r)$  if  $p, r \in S_3$ , and
- $V(S_1 \cup p, S_2, S_3 \setminus p) \geq V(S_1 \cup r, S_2, S_3 \setminus r)$  if  $p, r \in S_3$ .


$$(S_1 \cup p, S_2 \setminus p, S_3)$$


$$(S_1, S_2 \cup p, S_3 \setminus p)$$


$$(S_1 \cup p, S_2, S_3 \setminus p)$$



## A weighted game non being $I$ -complete

### Example

Consider the (3, 2) simple game with set of voters  $N = \{a, b, c\}$ :

$$W^m = \{(\emptyset, bc, a), (a, c, b)\} \quad \mapsto \quad L^M = \{(a, b, c), (c, a, b)\}$$

It is **weighted**

$$q = 0; \quad w(a) = (1, 0, 0), \quad w(b) = (0, 0, -1), \quad w(c) = (0, 0, -2).$$

but not  $I$ -complete:  $a \not\prec_I b$  and  $b \not\prec_I a$ , thus the game is not  $I$ -complete.



## Weighted and $I$ -complete $(3, 2)$ games

$I$ -complete  $(3, 2)$  games

weighted  
 $(3, 2)$  games

strongly weighted  $(3, 2)$  games

$SW \subseteq W \cap IC$




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(i)  $D^+$ -desirability.

$$p \succsim_{D^+} r \iff V(S_1 \cup p, S_2 \setminus p, S_3) \geq V(S_1 \cup r, S_2 \setminus r, S_3)$$

  
 $(S_1 \cup p, S_2 \setminus p, S_3)$


(ii)  $D^-$ -desirability.

$$p \succsim_{D^-} r \iff V(S_1, S_2 \cup p, S_3 \setminus p) \geq V(S_1, S_2 \cup r, S_3 \setminus r)$$

  
 $(S_1, S_2 \cup p, S_3 \setminus p)$

(iii)  $D^\pm$ -desirability.

$$p \succsim_{D^\pm} r \iff V(S_1 \cup p, S_2, S_3 \setminus p) \geq V(S_1 \cup r, S_2, S_3 \setminus r)$$

  
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- (i) A (3, 2) game is  $D^+$ -complete if either  $p \succsim_{D^+} r$  or  $r \succsim_{D^+} p$  for all pair  $p, r$  of voters.
- (ii) A (3, 2) game is  $D^-$ -complete if either  $p \succsim_{D^-} r$  or  $r \succsim_{D^-} p$  for all pair  $p, r$  of voters.
- (iii) A (3, 2) game is  $D^\pm$ -complete if either  $p \succsim_{D^\pm} r$  or  $r \succsim_{D^\pm} p$  for all pair  $p, r$  of voters.
- (iv) A (3, 2) game is complete if it is  $D^+$ -complete,  $D^-$ -complete and  $D^\pm$ -complete.



## Link between weights and desirability relations

Given two arbitrary players  $p$  and  $r$  in a weighted (3, 2) game, for any weight function representing it, we have:

- (i)  $w^+(p) \geq w^+(r) \Rightarrow p \succsim_{D^+} r$ .
- (ii)  $-w^-(p) \geq -w^-(r) \Rightarrow p \succsim_{D^-} r$ .
- (iii)  $w^+(p) - w^-(p) \geq w^+(r) - w^-(r) \Rightarrow p \succsim_{D^\pm} r$ .

### Corollary

*Every weighted (3, 2) game is a complete (3, 2) simple game for which the three relations  $\succsim_{D^+}$ ,  $\succsim_{D^-}$  and  $\succsim_{D^\pm}$  are transitive.*

### Corollary

*A strongly weighted (3, 2) game is an l-complete (3, 2) game.*

## Inclusion of types of VRA

$l$ -complete game  $\subset$  complete game

$\cup$

$\cup$

strongly weighted game  $\subset$  weighted game



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$l$ -complete game  $\subset$  complete game

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strongly weighted game  $\subset$  weighted game

$l$ -complete game  $\not\subseteq$  weighted game

weighted game  $\not\subseteq$   $l$ -complete game



## A small classification

Every  $(3, 2)$  simple game with:

- $n = 2$  is  $l$ -complete and moreover strongly weighted.
- $n = 3$  is complete. (\*)
- (\*) Although some of them are neither weighted nor  $l$ -complete.
- If  $n > 3$  one easily may find examples of  $(3, 2)$  simple games non-being complete.



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Banzhaf's (3,2) extension,  $(Bz_p[W])$ . Felsenthal and Machover introduced it in 1997. They consider the raw index as:

$$\eta_p[W] = |\{S : S \in W, S_{\downarrow p} \notin W\}|$$

The number  $\eta_p[W]$  counts the number of winning tripartitions in which  $p$  is positively decisive descending one single level of approval.

$$|\{S : p \in S_2, S \in W, S_{\downarrow p} \notin W\}| = |\{S : p \in S_1, S_{\downarrow p} \in W, S_{\downarrow\downarrow p} \notin W\}|$$



**Banzhaf's (3,2) extension**,  $(B_{z_p}[W])$ . Felsenthal and Machover introduced it in 1997. They consider the raw index as:

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$$\eta_p[W] = |\{S : S \in W, S_{\downarrow\downarrow p} \notin W\}|.$$

$B_{z_p}[W]$  is a numerical measure of decisiveness.

$$B_{z_p}[W] = \frac{\eta_p[W]}{\text{total number of tripartitions with } p \in S_1} = \frac{\eta_p[W]}{3^{n-1}}$$



We extend to weighted (3,2) games the idea by Brams and Affuso, which consists of using **generating functions** to obtain  $\eta[W]$  as the sum of some of its coefficients.

We build the *global generating function*  $f(x)$  and the *individual generating functions*  $f_i(x)$  in the following way:

$$f(x) = \prod_{p \in N} (x^{w^N(p)} + 1 + x^{w^Y(p)}) = \sum_{j=\bar{W}^N}^{\bar{W}^Y} \alpha_j x^j$$

$$f_i(x) = \prod_{p \neq i} (x^{w^N(p)} + 1 + x^{w^Y(p)}) = \sum_{j=\bar{W}^N - w^N(i)}^{\bar{W}^Y - w^Y(i)} \alpha_j x^j$$

where  $\bar{W}^Y = \sum_{p \in N} w^Y(p)$ ,  $\bar{W}^N = \sum_{p \in N} w^N(p)$ .



$$W^m = \{(12, \emptyset, 3), (1, 23, \emptyset), (23, 1, \emptyset)\} \equiv [2; (2, -1), (1, -2), (1, -1)]$$

$$f_1(x) = (1 + x + x^{-2})(1 + x + x^{-1}) = x^{-3} + x^{-2} + 2x^{-1} + 2 + 2x + x^2$$

$$f_2(x) = (1 + x^2 + x^{-1})(1 + x + x^{-1}) = x^{-2} + 2x^{-1} + 2 + 2x + x^2 + x^3$$

$$f_3(x) = (1 + x^2 + x^{-1})(1 + x + x^{-2}) = x^{-3} + x^{-2} + x^{-1} + 3 + x + x^2 + x^3$$

$$\left. \begin{array}{l} q - w^Y(1) = 0 \\ q - 1 = 1 \end{array} \right\} \left. \begin{array}{l} q = 2 \\ q - w^N(1) - 1 = 2 \end{array} \right\} \rightarrow \eta_1[W] = (\alpha_0 + \alpha_1) + a_2 = 4 + 1 = 5$$

$$\left. \begin{array}{l} q - w^Y(2) = 1 \\ q - 1 = 1 \end{array} \right\} \left. \begin{array}{l} q = 2 \\ q - w^N(2) - 1 = 3 \end{array} \right\} \rightarrow \eta_2[W] = \alpha_1 + (\alpha_2 + \alpha_3) = 2 + 2 = 4$$

$$\left. \begin{array}{l} q - w^Y(3) = 1 \\ q - 1 = 1 \end{array} \right\} \left. \begin{array}{l} q = 2 \\ q - w^N(3) - 1 = 2 \end{array} \right\} \rightarrow \eta_3[W] = \alpha_1 + \alpha_2 = 1 + 1 = 2$$

$$Bz[W] = \frac{1}{9}(5, 4, 2)$$



## Example

The UNSC voting system (as a simple game)

$$[39; 7, 7, 7, 7, 7, 1, 1, 1, 1, 1, 1, 1, 1, 1, 1]$$

The UNSC voting system (as a (3,2) game)

$$[ 9; \underbrace{(1, -6), \dots, (1, -6)}_5, \underbrace{(1, 0), \dots, (1, 0)}_{10} ]$$

Indices	Simple game		(3,2) game	
	Permanent M.	Nonpermanent M.	Permanent M.	Nonpermanent M.
$Rae(\mathcal{W})$	0, 5259	0, 5026	0, 3409	0, 337
$Bz(\mathcal{W})$	0, 0518	0, 0051	0, 0227	0, 0111
$Col^P(\mathcal{W})$	1	0, 0991	1	0, 3621
$Col^I(\mathcal{W})$	0, 0266	0, 0026	0, 0125	0, 0075
$KB(\mathcal{W})$	1	0, 5495	0, 738	0, 5747

Proportion of power for simple games ( $Bz \simeq 10 : 1$ )

Proportion of power for (3,2) games ( $Bz \simeq 2 : 1$ )

## A different notion of power?

In our notation we have  $\eta^+(W) + \eta^-(W) = \eta^\pm(W)$  while both  $\eta^+(W)$  and  $\eta^-(W)$  are independent.

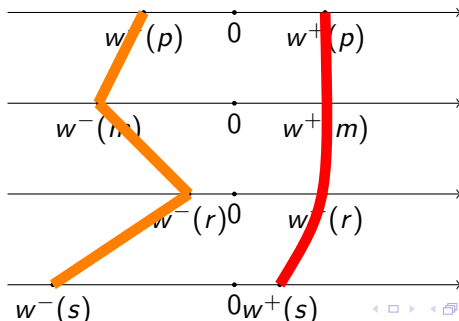




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We do think that Banzhaf power in the (3, 2) context should be conceived as **2-component power** rather than just as a single numerical measure.



Decomposition in 2-component power for the UNSC:

$$\begin{array}{lll} Bz_p^{YA}[V] = 0.0146 & Bz_p^{AN}[V] = 0.0080 & Bz_p^{YN}[V] = 0.02265 \\ Bz_p^{YA}[V] = 0.0111 & Bz_p^{AN}[V] = 0 & Bz_p^{YN}[V] = 0.01111 \end{array}$$



# THANKS FOR YOUR ATTENTION

